## CSCI 334: Principles of Programming Languages

Lecture 9: Lisp

Instructor: Dan Barowy Williams

#### Announcements

Midterm exam *next* Thursday

Optional Lisp homework posted later today

Feedback: slides insufficient (clarification?)

#### Choosing a language

Many good technical reasons

Few people decide for technical reasons

Seibel: Why do people get so religious about their computer languages?

**Bloch:** I don't know. But when you choose a language, you're choosing more than a set of technical trade-offs—you're choosing a community. It's like choosing a bar. Yes, you want to go to a bar that serves good drinks, but that's not the most important thing. It's who hangs out there and what they talk about. And that's the way you choose computer languages. Over time the community builds up around the language—not only the people, but the software artifacts: tools, libraries, and so forth. That's one of the reasons that sometimes languages that are, on paper, better than other languages don't win—because they just haven't built the right communities around themselves.



John McCarthy





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04004	-2 00001 4 0	04017 P6	TNX P5,K,1	
04005	0 76500 0 0	0043 P3	LRS 35	
04008	0 26000 0 0	4046	FMP X	
-04007	0 30000 1 0	4022	FAD A+1,J	
04010	1 77777 1 0	4011	TXI Pl,J,-1	
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429	CALL INFOL	0125
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433	WRITE UDIPUT TAPE OF TAPE TAPE TAPE TAPE	0127
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FORTRAN (circa 1956) (From <u>NASA Technical Note D-1737</u>)

LISP (circa 1958)

LISP is a "functional" language

- programs are modeled after math. functions
- no statements, only expressions
- no "mutable" variables, only declarations
- therefore, the effect of running a program ("evaluation") is purely the effect of applying a function to an input.

What does this mean?

- there is no hidden state (globals in C) or statechange rules (as in Java).
- complex parts are composed out of simple parts
- we care about what a program means and not what it does: you can understand a program without having to mentally execute it
- programs are easier to reason about

functional language (again)

- programs are modeled after math. functions
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LISP is inspired by the lambda calculus

• everything either a value or a funct. of a value value: 1

function of two values: (+ 1 1)

- syntax is (mind-numbingly) regular
   functions: (function-name arguments ...)
   values: anything that is not a function
- evaluating a function produces a value:

**(**+ 1 1**)** → 2

Statements vs. expressions

• A *statement* is an operation that changes the state of the computer

Java: i++

value stored at location i incremented by one

An *expression* is a combination of functions and values that yields a new value
 Lisp: (+ i 1)
 evaluating + with i and 1 returns i + 1

REPL (read-eval-print loop)

Batch mode

### Mutable variables

• If you can update a variable, your language has mutable variables

Java: int i = 3;

- i = 4;
- Notice that both lines of code are statements
- (pure) Lisp does not have mutable variables



Lisp: (defun my-func (i) ...)

(my-func 3)

or the shorter

((lambda (i) ...) 3)

- Notice that all of the above are expressions
- In fact, functions are the only way to bind values to names in (pure) Lisp







Lisp syntax: car and cdr

• Access the first element of a cons cell with car

(car (cons 1 2)) = 1

• Access the second element with cdr

```
(cdr (cons 1 2)) = 2
```

• What's the value of the following expression?

```
(car '(1 2 3))
```

• What about this?

```
(cdr '(1 2 3))
```

### Lisp syntax: functions

- Everything else is a function (or special form)
- There are a bunch of built-in functions

(car ...)

(cdr ...)

(append ...), etc.

 And you can define your own (defun my-func (arg) (value))







```
C:
When you want to use a variable, you have to
allocate it first, then decallocate it when done.
MyObject *m = malloc(sizeof(MyObject));
m->foo = 2;
m->bar = 3;
... do stuff with m ...
free(m);
```

Memory management

# Memory management

• Lisp and Java:

```
You barely need to think about this at all.
MyObject m = new MyObject(2,3);
... do stuff with m ...
(cons 2 3)
```