

CSCI 334:
Principles of Programming Languages

Lecture 13: Parsing

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Announcements

HW7 Q5 updated (it's a tad easier now)

You can find the Parser lib linked in the reading.

HW4 resubmission: last day to submit is today

HW8 is your project proposal: think about who you want to work with. I am happy to help find partners.

Compound Types:

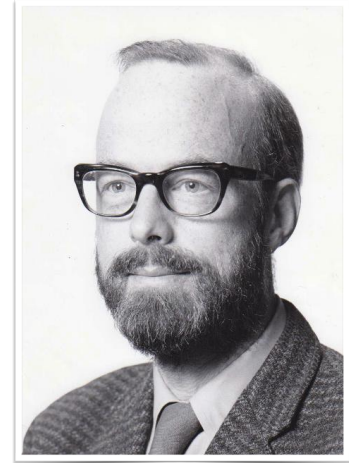
Records, 
Lists, 
Tuples, 
ADTs

Algebraic Data Types

not
Abstract Data Types

Hoare Property

"There are two ways of constructing a software design: One way is to make it so simple that there are obviously no deficiencies. and the other way is to make it so complicated that there are no obvious deficiencies." — C.A.R. Hoare



Hoare Property

ADTs make the structure of a program's logic
"more obvious."

ADT (Java)

```
public static final int NORTH = 1;
public static final int SOUTH = 2;
public static final int EAST = 3;
public static final int WEST = 4;

public move(int x, int y, int dir) {
    switch (dir) {
        case NORTH: ...
        case ...
    }
}
```

ADT (F#): "discriminated unions"

```
type Direction =  
    North | South | East | West;  
  
let move x y dir =  
    match x,y,dir with  
    | x,y,North -> x,y-1  
    | x,y,South -> x,y+1
```

- Above is an "incomplete pattern"
- ML will warn you when you've missed a case!
 - "proof by exhaustion"

ADTs can have parameters

```
type Shape =  
    | Rectangle of float * float  
    | Circle of float
```

- Pattern match to extract parameters

```
let s = Rectangle(1.0,4.0)  
match s with  
| Rectangle(w,h) -> ...  
| Circle(r) -> ...
```

ADTs can have named parameters

```
type Shape =  
    | Rectangle of width: float * height: float  
    | Circle of radius: float
```

- Names are useful for initialization and documentation.

```
let s = Rectangle(height = 1.0, width = 4.0)
```

ADTs can even be recursive and generic

```
type MyList<'a> =  
    | Empty  
    | NonEmpty of head: 'a * tail: MyList<'a>
```

```
> NonEmpty(2, Empty);;  
val it : MyList<int> = NonEmpty (2,Empty)
```

handling errors with ADTs and patterns

- Another example: handling errors.
- F# has exceptions (like Java)
- But an alternative, **easy** way to handle many errors is to use the option type:

```
type Option<'a> = None | Some of 'a
```

handling undefinedness with patterns

Write a function `get_nth` that takes a **list of strings** and **an int n** and **returns the nth element** of the list, where the head is 1st.

handling undefinedness with patterns

```
let rec get_nth xs n =  
  match xs, n with  
  | [], _ -> None  
  | (x::xs), 1 -> Some x  
  | (x::xs), n ->  
    if n > 1 then get_nth xs (n-1) else None
```

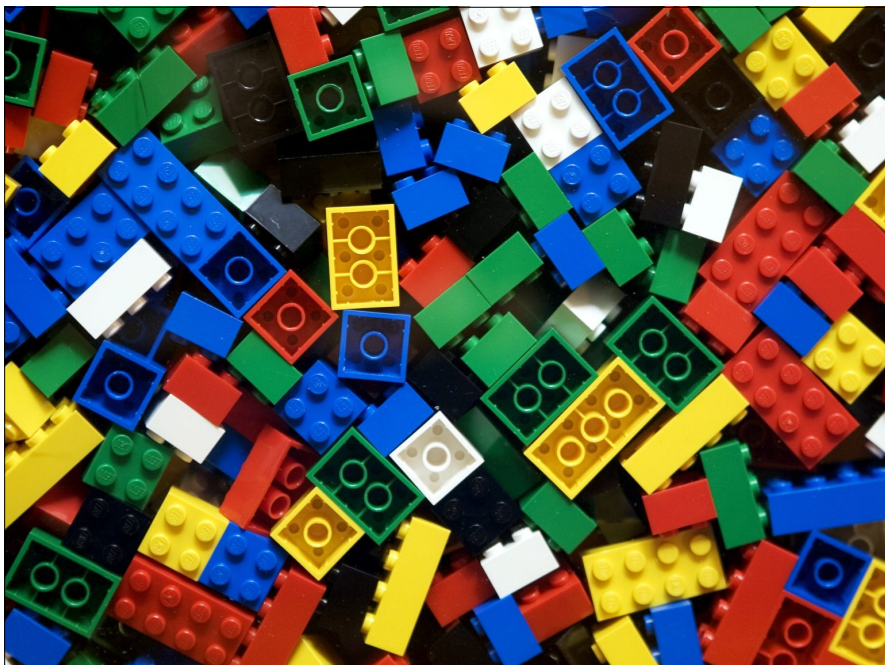
option type

- Why option?
- option is a **data type**;
not handling errors is a **static type error**!

handling errors with patterns

```
> get_nth [1;2;3;4] 3;;  
val it : int option = Some 3  
> get_nth [1;2;3;4] 0;;  
val it : int option = None  
> get_nth [1;2;3;4] 5;;  
val it : int option = None  
> get_nth [1;2;3;4] -2;;  
val it : int option = None
```

Parser Combinators



Parser Combinators

- A kind of **recursive decent** parser.
- Recursive descent parser:
A **top-down** parser built from a set of **mutually recursive procedures** where each such procedure **usually implements one of the productions** of the grammar.

Basic Primitives

- Input
`type Input = string * bool`
- Output
`type Outcome<'a> =
| Success of result: 'a * remaining: Input
| Failure`

Basic Primitives

- A parser is
`type Parser<'a> = Input -> Outcome<'a>`
- Keep in mind: a parser *is a function*.

Two varieties of parser

- Parsers that **consume input**. Correspond with grammar terminals.
- Parsers that **combine parsers**. Correspond with grammar non-terminals.
- For flexibility, you can also have **parsers that do both**.

A very simple terminal parser

- To parse a given char

```
pchar(c: char) : Parser<char>
```

- Notice that the **generic type** inside <brackets> is the **return type of the parser**.
- So `pchar` returns a `char`.

How to use it

- `(pchar 'z') input`
- input must be "prepared" first.
- ```
> let input = "zoo";;
 val input : string = "zoo"
> let i = prepare input;;
 val i : Input = ("zoo", true)
> (pchar 'z') i;;
 val it : Outcome<char> = Success ('z', ("oo", true))
```

## A very simple combining parser

- To parse two things in sequence:

```
pseq : p1:Parser<'a> -> p2:Parser<'b> ->
f:('a * 'b -> 'c) -> Parser<'c>
```

- It looks more complicated than it is.
- Let's look at each part.

## A very simple combining parser

- `pseq :`  

```
 p1:Parser<'a>
 ->
 p2:Parser<'b>
 ->
 f:('a * 'b -> 'c) -> Parser<'c>
```
- `p1` is a parser.

## A very simple combining parser

- pseq :  
    p1:Parser<'a>  
    ->  
    p2:Parser<'b>  
    ->  
    f:('a \* 'b -> 'c) -> Parser<'c>
- p2 is a parser.

## A very simple combining parser

- pseq :  
    p1:Parser<'a>  
    ->  
    p2:Parser<'b>  
    ->  
    f:('a \* 'b -> 'c) -> Parser<'c>
- f is a function that takes the result of p1 and p2 and does something with it. That something is up to **you**.

## How to use it

- pseq (pchar 'z') (pchar 'o') id
- id is F#'s identity function.
- Let's play with this in fsharp.i.

## More details

- It is **critical** that you read the "Parser Combinators" reading.
- I suggest that you **sit down, uninterrupted, for an hour or two**, and **work through the examples** in fsharp.i.
- The reading builds the `Parsers.fs` library that you are given for HW7.



## Example: brace language

- An *expression* is a sequence of *terms*, consisting of *at least one term*.
- A *term* is either 'aaa', 'bbb', or a *brace expression*.
- A *brace expression* is '{', followed by an *expression*, followed by '}'.